

## **Irreducible Surfaces**

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Supplementary material for

**“A dynamic programming algorithm for  
RNA structure prediction including pseudoknots”**

**by Elena Rivas and Sean R. Eddy.**

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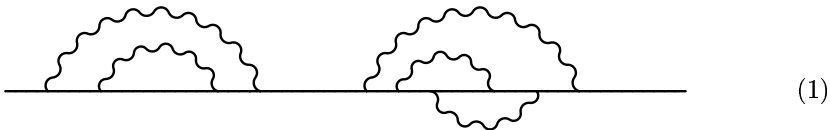
In this supplementary material we are give a more precise mathematical definition and some properties of Irreducible Surfaces (ISs).

Irreducible Surfaces are of importance because they constitute the basic elements of RNA folding: hairpin loops, stems, bulges, internal loops, and multiloops. ISs can also be pseudoknotted. In this regard, ISs are an extension of the “k-loops” introduced by Sankoff (1985) which are inherently nested.

Our algorithm is an expansion in ISs, that is, it is an attempt to express a given full structure in terms of its basic constituents, much in the same way as Schwinger-Dyson diagrams do for quantum field theory. In this supplementary material we introduce some relevant concepts about the topology of the graphs we are using.

**Definition 1 (Graph)** A graph is any continuous fragment of RNA that folds into an arbitrary non-trivial secondary structure.

Example of a graph,

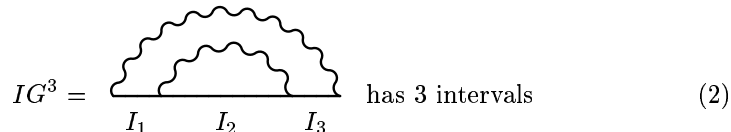


**Definition 2 (Irreducible Graph)** A graph is irreducible when the RNA backbone cannot be disrupted, without disrupting any secondary interactions. The above graph is reducible to two irreducible components.

**Definition 3 (Interval)** We define an interval,  $I$ , of an irreducible graph (IG) as set of points in between two secondary interactions, including the paired vertices.

**Definition 4 (Order of an IG)** If a given IG has  $\gamma$  intervals, the graph is said to be of order  $\gamma$ ,  $IG^\gamma$ .

Example, the graph in (1) contains two irreducible graphs,





**Theorem 2** A surface is irreducible if and only if the intervals that constitute the surface are disjoint.

*Proof:*

$\Leftarrow$  If the intervals are disjoint, no H-bond is shared by any two intervals, therefore if an H-bond is broken, the surface becomes an open path.

$\Rightarrow$  Suppose a surface  $S^\gamma$  in which two intervals  $I_i, I_j$  are not disjoint. Then, intervals  $I_i$  and  $I_j$  share an H-bond. Because the surface closed in itself by definition, there is another pair of non-disjoint intervals  $I'_i$  and  $I'_j$  that share the other end of the H-bond. Now break that common interaction, then  $S^{\gamma-2}(I_1, \dots, I_i \cup I_j, \dots, I'_i \cup I'_j, \dots, I_\gamma)$  is a surface.

**Corollary 1** A surface is irreducible if and only if no H-bond has to be crossed more than one time in traversing the surface.

*Proof:*

The intervals are disjoint if and only if no H-bond has to be crossed more than once.

**Property 2** A surface of order  $\gamma$  always contains at least one irreducible surface  $IS^\beta$ , such that  $\beta \leq \gamma$ .

*Proof:*

Suppose that  $S^\gamma$  has  $\beta \leq \gamma$  H-bonds  $\Rightarrow \exists \beta$  disjoint intervals that form a surface  $\Rightarrow$  there is an  $IS^\beta \subset S^\gamma$ . This  $IS^\beta$  is the one of higher order that a  $IG^\gamma$  can have.

In summary, our algorithm is a recursion for generic graphs (the  $wx, vx, vhx, whx, \dots$  matrices) in terms of irreducible surfaces. Properties 1 and 2 assure us that for a given generic graph there is a recursion possible in terms of ISs, and that the recursion has a finite number of terms.